## Q. Design Turing Machine for language $\mathrm{L}=\left\{0^{\mathrm{n}} \mathbf{1}^{\mathrm{n}} \mathbf{2}^{\mathrm{n}} \mid\right.$ $n \geq 1\}$

## Solution:

Before designing the required Turing machine $M$, let us evolve a procedure for processing the input string 112233. After processing, we require the ID to be of the form $b b b b b b q_{7}$. The processing is done by using five steps:

Step $1 q_{1}$ is the initial state. The R/W head scans the leftmost 1 , replaces 1 by $b$, and moves to the right. $M$ enters $q_{2}$.
Step 2 On scanning the leftmost 2, the R/W head replaces 2 by $b$ and moves to the right. $M$ enters $q_{3}$.

Step 3 On scanning the leftmost 3, the R/W head replaces 3 by $b$, and moves to the right. $M$ enters $q_{4}$.

Step 4 After scanning the rightmost 3, the R/W heads moves to the left until it finds the leftmost 1 . As a result. the leftmost 1,2 and 3 are replaced by $b$.

Step 5 Steps 1-4 are repeated until all 1's, 2's and 3's are replaced by blanks. The change of IDs due to processing of 112233 is given as

$$
\begin{gathered}
q_{1} 112233\left|-b q_{2} 12233\right|-b 1 q_{2} 2233\left|-b 1 b q_{3} 233\right|-b 1 b 2 q_{3} 33 \\
\vdash b 1 b 2 b q_{4} 3-b 1 b_{2} q_{5} b 3-b 1 b q_{5} 2 b 3\left|-b 1 q_{5} b 2 b 3\right|-b q_{5} 1 b 2 b 3 \\
\vdash q_{6} b 1 b 2 b 3\left|-b q_{1} 1 b 2 b 3\right|-b b q_{2} b 2 b 3 \mid-b b b q_{2} 2 b 3 \\
\vdash b b b b q_{3} b 3-b b b b b q_{3} 3-b b b b b b q_{4} b \vdash b b b b b q_{7} b b
\end{gathered}
$$

Thus.

$$
q_{1} 112233 \vdash q_{7} b b b b b b
$$

As $q_{7}$ is an accepting state, the input string 112233 is accepted.

Now we can construct the transition table for $M$.

| Present state | Input tape symbol |  |  |  |
| :---: | :---: | :---: | :---: | :---: |
|  | 1 | 2 | 3 | $b$ |
| $\rightarrow q_{1}$ | $b R q_{2}$ |  |  | $b R q_{1}$ |
| $q_{2}$ | $1 R q_{2}$ | $b R q_{3}$ | $b R q_{2}$ |  |
| $q_{3}$ |  | $2 R q_{3}$ | $b R q_{4}$ | $b R q_{3}$ |
| $q_{4}$ |  |  | $3 L q_{5}$ | $b L q_{7}$ |
| $q_{5}$ | $1 L q_{6}$ | $2 L q_{5}$ |  | $b L q_{5}$ |
| $q_{5}$ | $1 L q_{6}$ |  | $b R q_{1}$ |  |
| $q_{7}$ |  |  |  |  |

It can be seen from the table that strings other than those of the form $0^{n} 1^{n} 2^{n}$ are not accepted. It is advisable to compute the computation sequence for strings like $1223,1123,1233$ and then see that these strings are rejected by $M$.

